

Adam Paweł Wojda

## EXTREMAL TRACEABLE GRAPHS WITH NON-TRACEABLE EDGES

**Abstract.** By  $\text{NT}(n)$  we denote the set of graphs of order  $n$  which are traceable but have non-traceable edges, i.e. edges which are not contained in any hamiltonian path. The class  $\text{NT}(n)$  has been considered by Balińska and co-authors in a paper published in 2003, where it was proved that the maximum size  $t_{\max}(n)$  of a graph in  $\text{NT}(n)$  is at least  $(n^2 - 5n + 14)/2$  (for  $n \geq 12$ ). The authors also found  $t_{\max}(n)$  for  $5 \leq n \leq 11$ .

We prove that, for  $n \geq 5$ ,  $t_{\max}(n) = \max\{\binom{n-2}{2} + 4, \binom{n - \lfloor \frac{n-1}{2} \rfloor}{2} + \lfloor \frac{n-1}{2} \rfloor^2\}$  and, moreover, we characterize the extremal graphs (in fact we prove that these graphs are exactly those already described in the paper by Balińska *et al.*). We also prove that a traceable graph of order  $n \geq 5$  may have at most  $\lceil \frac{n-3}{2} \rceil \lfloor \frac{n-3}{2} \rfloor$  non traceable edges (this result was conjectured in the mentioned paper by Balińska and co-authors).

**Keywords:** traceable graph, non-traceable edge.

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### 1. INTRODUCTION

Throughout the paper we use the standard terminology, unless stated otherwise. An edge  $e$  of a graph  $G$  is called *traceable* if it is contained in a hamiltonian path of  $G$ , otherwise it is called *non-traceable* edge of  $G$  (*NT-edge*). If  $G$  has a hamiltonian path and at least one non-traceable edge, then  $G$  is called *NT-graph*. The set of all NT-graphs of order  $n$  is denoted by  $\text{NT}(n)$ . NT-graphs were introduced and studied in [1].

It is easy to check that the least number  $n$  for which  $\text{NT}(n)$  is not empty is  $n = 5$  and, moreover,  $\text{NT}(n)$  is not empty for every  $n \geq 5$ .

By  $t_{\max}(n)$  we denote the maximum size of NT-graph of order  $n$ . In [1] it was proved that

$$t_{\max}(n) = (3n^2 - 4n + 1)/8 \quad \text{for } n = 7, 9, \quad (1)$$

$$t_{\max}(n) = (n^2 - 5n + 14)/2 \quad \text{for } n = 5, 6, 8, 10, 11, \quad (2)$$

$$t_{\max}(n) \geq (n^2 - 5n + 14)/2 \quad \text{for } n \geq 12. \quad (3)$$

Balińska *et al.* [1] have also indicated that

$$\overline{K}_2 * K_2 * K_{n-4} \quad \text{for } n = 8 \quad \text{and for } n \geq 10, \quad (4)$$

$$\overline{K}_{\frac{n-1}{2}} * K_{\frac{n-1}{2}} * K_1 \quad \text{for } n = 5, 7, 9, 11, \quad (5)$$

$$\overline{K}_{\frac{n-2}{2}} * K_{\frac{n-2}{2}} * K_2 \quad \text{for } n = 6 \quad \text{and for } n = 8, \quad (6)$$

(where  $*$  denotes the join of graphs) are NT-graphs with size  $(3n^2 - 4n + 1)/8$  for  $n = 7, 9$  and  $(n^2 - 5n + 14)/2$  for  $n \geq 5, n \neq 7, 9$ .

We prove that the bounds found in [1] are in fact the best possible and, moreover, the family of graphs described by (4)–(6) is the set of all NT-graphs of order  $n$  and maximum size.

**Theorem 1.** *The maximum size  $t_{\max}(n)$  of NT-graph of order  $n$  is given by the formula*

$$t_{\max}(n) = \max \left\{ \binom{n-2}{2} + 4, \binom{n - \lfloor \frac{n-1}{2} \rfloor}{2} + \left\lfloor \frac{n-1}{2} \right\rfloor^2 \right\}.$$

Moreover, the NT-graphs of size  $t_{\max}(n)$  are the following.

1.  $\overline{K}_2 * K_2 * K_{n-4}$  for  $n = 8$  and for  $n \geq 10$ ,
2.  $\overline{K}_{\frac{n-1}{2}} * K_{\frac{n-1}{2}} * K_1$  for  $n = 5, 7, 9, 11$ ,
3.  $\overline{K}_{\frac{n-2}{2}} * K_{\frac{n-2}{2}} * K_2$  for  $n = 6, 8$ .

Note that for  $n = 8$  and for  $n = 11$  there are two extremal graphs. An easy computation shows that  $\binom{n-2}{2} + 4 \geq \binom{n - \lfloor \frac{n-1}{2} \rfloor}{2} + \lfloor \frac{n-1}{2} \rfloor^2$  for  $n = 5, 6, 8$  and for  $n \geq 10$ .

For every integer  $n \geq 5$  the graph  $\overline{K}_2 * K_2 * K_{n-4}$  contains exactly one NT-edge. It is also very easy to check that for  $n \geq 5$  the graph  $\overline{K}_{\lfloor \frac{n-1}{2} \rfloor} * K_{\lfloor \frac{n-1}{2} \rfloor} * K_{n-2\lfloor \frac{n-1}{2} \rfloor}$  has exactly  $\binom{\lfloor \frac{n-1}{2} \rfloor}{2}$  NT-edges. In [1], for every  $n \geq 5$ , Balińska *et al.* presented a graph of order  $n$  with  $\lceil \frac{n-3}{2} \rceil \lfloor \frac{n-3}{2} \rfloor$  NT-edges. They also conjectured that the maximum number of NT-edges of a graph in  $\text{NT}(n)$  is equal to  $\lceil \frac{n-3}{2} \rceil \lfloor \frac{n-3}{2} \rfloor$ . We prove this conjecture.

**Theorem 2.** *Let  $n \geq 5$  and let  $b_{\max}(n)$  denote the maximum number of NT-edges of a graph in  $\text{NT}(n)$ . Then*

$$b_{\max}(n) = \left\lceil \frac{n-3}{2} \right\rceil \left\lfloor \frac{n-3}{2} \right\rfloor.$$

## 2. PROOF OF THEOREM 1

The proof of Theorem 1 follows from the two below quoted theorems in [2] (these theorems are in fact corollaries from Theorem 3 of [2], which is a much more general result).

**Theorem 3** (see [2]). *Let  $G$  be a graph of order  $n \geq 5$  and minimum vertex degree  $\delta(G) \geq 2$ . If  $G$  is a non-hamiltonian graph of maximum size then either*

$$G = \overline{K}_2 * K_2 * K_{n-4}$$

or  $n \in \{5, 6, 7, 8, 9, 11\}$  and

$$G = \overline{K}_{\lfloor \frac{n-1}{2} \rfloor} * K_{\lfloor \frac{n-1}{2} \rfloor} * K_{n-2\lfloor \frac{n-1}{2} \rfloor}.$$

Note that the set of extremal graphs in Theorem 3 is exactly the same as that in Theorem 1.

**Theorem 4** (see [2]). *If  $G$  is a non-hamiltonian connected graph of order  $n \geq 4$  and of maximum size, then either*

$$G = K_1 * K_2 * K_{n-3}$$

or else  $n = 6$  and

$$G = \overline{K}_3 * K_3.$$

Let  $G$  be a graph of order  $n \geq 5$  and size

$$\|G\| \geq \max \left\{ \binom{n-2}{2} + 4, \binom{n - \lfloor \frac{n-1}{2} \rfloor}{2} + \left\lfloor \frac{n-1}{2} \right\rfloor^2 \right\}.$$

We shall prove that every edge of  $G$  is contained in a hamiltonian path of  $G$  or  $G$  is one of the extremal graphs listed in the assertion of Theorem 1.

Let us first suppose that there is a vertex  $v$  in  $G$  such that  $d_G(v) \leq 1$  (by  $d_G(x)$  we denote the degree of the vertex  $x$  in  $G$ ). Note that then  $d_G(v) = 1$  since  $G$  is traceable. Then  $\|G - \{v\}\| = \|G\| - 1$ . We now apply Theorem 4 and prove that  $G - \{v\}$  is hamiltonian connected and check that then every edge of  $G$  is contained in a hamiltonian path of  $G$ .

It is very easy to see that if a graph is hamiltonian then every of its edges lies in a hamiltonian path – this fact was also observed in [1].

Suppose now that  $\delta(G) \geq 2$ . By Theorem 3, either  $G$  is hamiltonian, and therefore contains every edge in a hamiltonian path, or it is one of extremal graphs of Theorem 1 and the proof is finished.

## 3. PROOF OF THEOREM 2

Since it has been shown in [1] that for every  $n \geq 5$  there is a graph of order  $n$  which is traceable and has  $\lfloor \frac{n-3}{2} \rfloor \lceil \frac{n-3}{2} \rceil$  non-traceable edges, it remains to prove that in every traceable graph  $G$  of order  $n \geq 5$  there are at most  $\lfloor \frac{n-3}{2} \rfloor \lceil \frac{n-3}{2} \rceil$  non-traceable edges.

Let  $(x_1, \dots, x_n)$  be a hamiltonian path of a traceable graph  $G$  of order  $n$ . The following facts are very easy to check.

- (i) every edge  $x_i x_{i+1}$  (for  $i = 1, \dots, n-1$ ) is traceable,
- (ii) every edge  $x_1 x_i$  ( $1 < i \leq n$ ) is traceable and every edge  $x_j x_n$  ( $1 \leq j < n$ ) is traceable,
- (iii) if  $i > 1, s \geq 2$  and  $i+s < n$ , then at most one of the edges  $x_i x_{i+s}$  and  $x_{i+1} x_{i+s+1}$  is non-traceable (if both edges  $x_i x_{i+s}$  and  $x_{i+1} x_{i+s+1}$  are present in  $G$ , then they are traceable).

By (i)–(ii) only the edges of the form  $x_i x_{i+s}$  with  $i > 1, 2 \leq s \leq n-3$  and  $i+s \leq n$  may be non traceable. Moreover, by (iii), for any fixed  $s, 2 \leq s \leq n-2$ , we may have at most  $\lceil \frac{n-s-2}{2} \rceil$  non-traceable edges of the form  $x_i x_{i+s}$ . Hence, the maximum number  $b(G)$  of non-traceable edges in  $G$  verifies

$$b(G) \leq \sum_{s=2}^{n-3} \left\lceil \frac{n-s-2}{2} \right\rceil.$$

One may now prove easily, for instance by induction, the formula

$$\sum_{s=2}^{n-3} \left\lceil \frac{n-s-2}{2} \right\rceil = \left\lfloor \frac{n-3}{2} \right\rfloor \left\lceil \frac{n-3}{2} \right\rceil$$

for  $n \geq 5$ .

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Adam Paweł Wojda  
wojda@uci.agh.edu.pl

AGH University of Science and Technology  
Faculty of Applied Mathematics  
Chair of Discrete Mathematics  
al. Mickiewicza 30, 30-059 Cracow, Poland

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